

Solution Methods for the Analysis of Large  
Markov Chains  
(CTMCs)

*Rashid Mehmood*

Swansea University

MATHEMATICS OF NETWORKS

(Sixth Meeting)

29 June 2007

# Context

---

- Stochastic Modelling
  - unpredictability is inherent in most real-life systems
    - e.g. communication networks
- Tools for analyzing such systems:
  - e.g. discrete-state models
  - e.g. Markov Chains, CTMCs (continuous time)
- Behaviour of physical systems can be described by:
  - a set of states and
  - state to state transitions

# Context . . .

---

- Three steps in performance evaluation:
  - 1. specification
    - Model description:  
e.g. queueing networks, stochastic Petri nets, . . .
    - Properties specification:  
probabilistic logics (CSL, PCTL, . . .)
  - 2. state space generation
    - matrix generation from the formalism
  - 3. computing performance measures
    - e.g. solving  $Ax = 0$  for steady state solution
- Problem: State-space explosion
  - we concentrate on phase three (*computationally expensive*):
    - numerical solution of  $Ax = 0$

# State-Space Explosion

---

- The task is the iterative solution of  $Ax = 0$ 
  - requires storage of the matrix  $A$  and the vector(s)  $x$
- Explicit techniques
  - disk-based storage of matrix (out-of-core)
  - distribution to parallel nodes
  - disk-based + parallel
- Implicit techniques (compact representation of matrix)
  - e.g. symbolic: BDD-based data structures
- The aim
  - extend the size of solvable models
  - develop new solution techniques
  - seek improvements in the existing
  - build realistic models of interesting applications

# Contents

---

- Numerical Computations and the available Methods
- Matrix Storage
  - Sparse and symbolic storage schemes
- The out-of-core solution Methods
- Parallel Symbolic implementation of Jacobi solution method
- Applications I have been looking at
- Summary and Future Work

# The Numerical Computations

---

- Steady state computations

- $\pi Q = 0, \sum_{i=0}^{n-1} \pi_i = 1$

- Direct Methods

- Modify the matrix
  - Problem: *fill-in*

- Iterative methods

- Based on matrix vector products (MVP)

# Iterative Methods for $Ax = b$

---

- *Residual*  $= b - Ax$ 
  - Make initial guess for the solution vector  $x$
  - Generate successive approximations
  - **Until** residual falls below some prescribed value
  - Matrix is unchanged
- **Basic Iterative methods**
  - Jacobi, Gauss-Seidel, SOR
  - Iteratively, modify components of solution vector
  - Convergence tested for residual or error vector
- **Projection methods**
  - Extract an approximate solution  $x$  from a subspace of  $\mathbb{R}^n$
  - Orthogonalisation, nonstationary, parameter-free
  - Krylov subspace methods

# Transition Matrices

---

- Continuous Time Markov Chains

- Transition rate matrix  $Q$

- $S \times S \rightarrow \mathbb{R}$

- Diagonal entries:  $q_{ii} = - \sum_{i \neq j} q_{ij}$

- Properties

- Very large, sparse

- Number of distinct values depends on the model



# Transition Matrices (Polling System)

---

- Case Study: Cyclic Server Polling System [Ibe and Trivedi, 1990]
  - $k$  stations or queues and a server
  - server polls the stations in a cycle looking for jobs

$k$	states ( $n$ )	off-diagonal nonzero ( $a$ )	$a/n$	MB per $\pi$
17	3,342,336	31,195,136	9.3	26
18	7,077,888	69,599,232	9.8	54
20	31,457,280	340,787,200	10.8	240
21	66,060,288	748,683,264	11.3	513
22	138,412,032	1,637,875,712	11.8	1,056
23	289,406,976	3,569,352,704	12.3	2,204
24	603,979,776	7,751,073,792	12.8	4,608
25	1,258,291,200	16,777,216,000	13.3	9,598

# Case Studies (Kanban and FMS Systems)

---

- Case Study: FMS [Ciardo and Tilgner , 1993]
  - Flexible Manufacturing System
  - models a system with three machines
  - machines process different types of parts
  - model parameter  $k$ : max. no. of parts machine can handle
  
- Case Study: Kanban [Ciardo and Tilgner , 1996]
  - Kanban Manufacturing System
  - models a system with four machines
  - parameter  $k$  is the max. no. of machine jobs at a time

# Sparse Matrix Storage Schemes

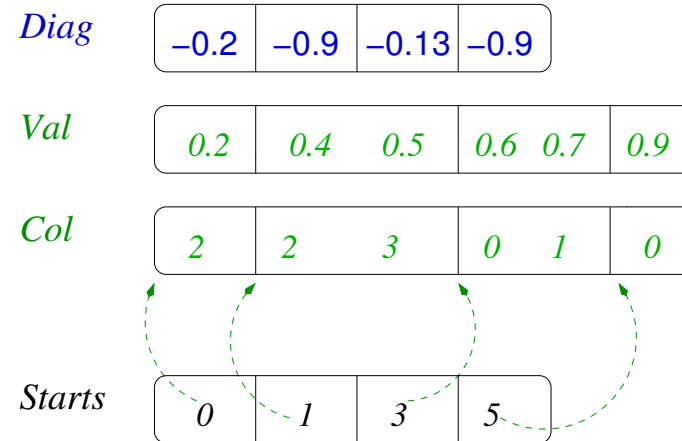
---

- Coordinate format, Compressed Sparse Row (CSR) schemes
  - Store whole matrix with no distinction between diagonal and off-diagonal entries
- Modified (compressed) Sparse Row (MSR)
  - Four arrays:
    - double `Diag`[ $n$ ]
    - double `Val`[ $a$ ]
    - int `Col`[ $a$ ]
    - int `Starts`[ $n$ ]
  - Bytes required:  $12(a + n)$

# Sparse Matrix Storage Schemes...

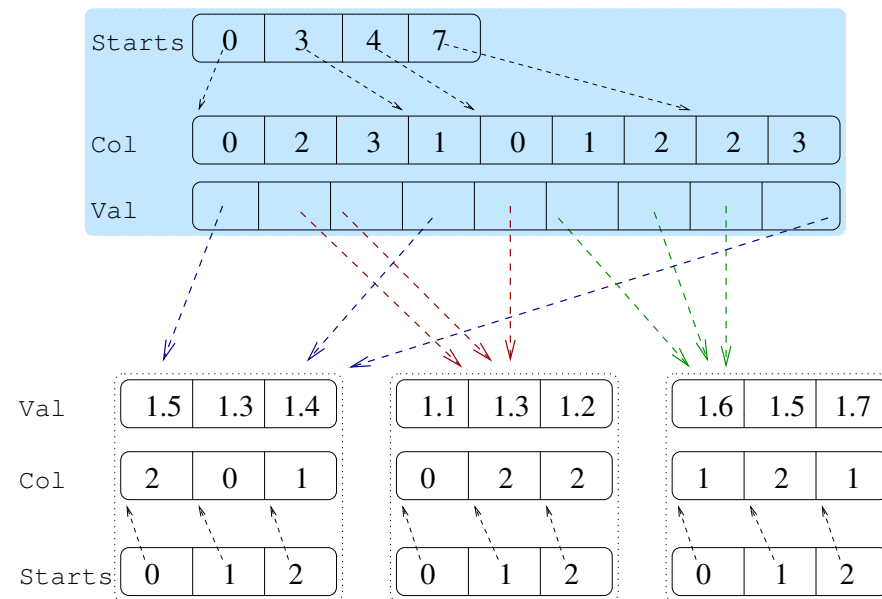
$$\begin{pmatrix} -0.2 & 0 & 0.2 & 0 \\ 0 & -0.9 & 0.4 & 0.5 \\ 0.6 & 0.7 & -0.13 & 0 \\ 0.9 & 0 & 0 & -0.9 \end{pmatrix}$$

$$n = 4, a = 6$$



# The Symbolic Representation of a CTMC

0 0 1.5		1.1 0 0	1.1 0 0
1.3 0 0		0 0 1.3	0 0 1.3
0 1.4 0		0 0 1.2	0 0 1.2
	0 0 1.5		
	1.3 0 0		
	0 1.4 0		
1.1 0 0	0 1.6 0	0 1.6 0	
0 0 1.3	0 0 1.5	0 0 1.5	
0 0 1.2	0 1.7 0	0 1.7 0	
		0 1.6 0	0 0 1.5
		0 0 1.5	1.3 0 0
		0 1.7 0	0 1.4 0

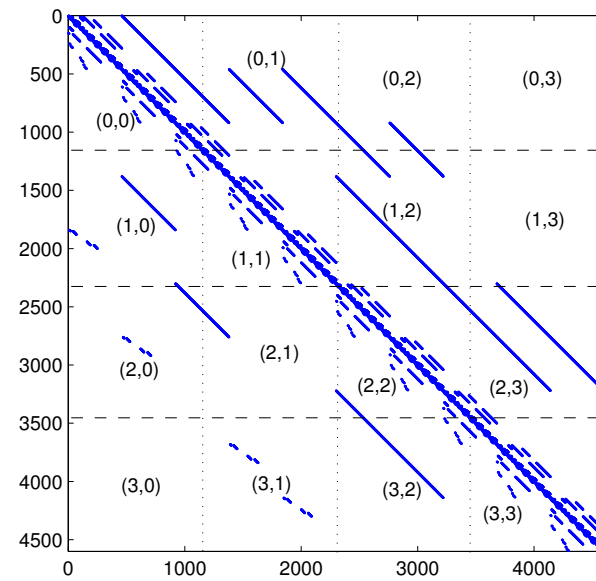
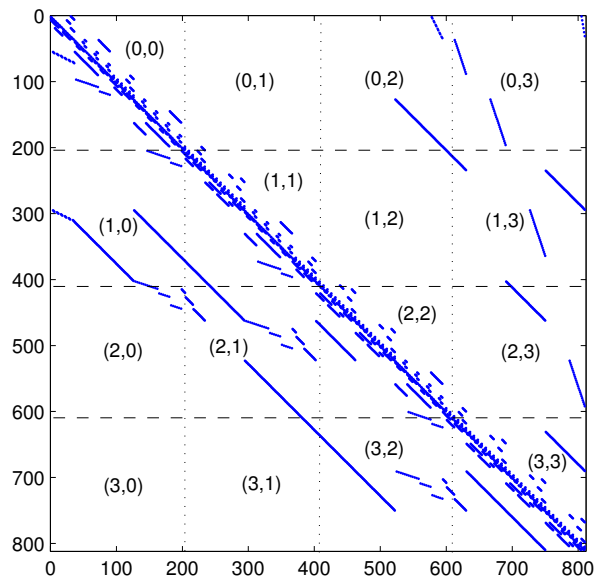


# Comparison of Storage Methods

$k$	States ( $n$ )	Off-diagonal nonzeros ( $a$ )	$a/n$	Memory for Matrix (MB)				Vector (MB)
				$MSR$	$Ind. MSR$	$Comp. MSR$	$MTBDDs$	
FMS models								
6	537,768	4,205,670	7.82	50	24	17	4	4
10	25,397,658	234,523,289	9.23	2,780	1,366	918	137	194
13	216,427,680	2,136,215,172	9.87	25,272	12,429	8,354	921	1,651
14	403,259,040	4,980,958,020	12.35	58,540	28,882	19,382	1,579	3,077
15	724,284,864	9,134,355,680	12.61	107,297	52,952	35,531	2,676	5,526
Kanban models								
4	454,475	3,979,850	8.76	47	23	16	1	3.5
6	11,261,376	115,708,992	10.27	1,367	674	452	6	86
9	384,392,800	4,474,555,800	11.64	52,673	25,881	17,435	99	2,933
10	1,005,927,208	12,032,229,352	11.96	141,535	69,858	46,854	199	7,675
Polling System								
15	737,280	6,144,000	8.3	73	35	24	1	6
21	66,060,288	748,683,264	11.3	8,820	4,334	2,910	66	504
24	603,979,776	7,751,073,792	12.8	91,008	44,813	30,067	144	1,136
25	1,258,291,200	16,777,216,000	13.3	196,800	96,960	65,040	317	1,190

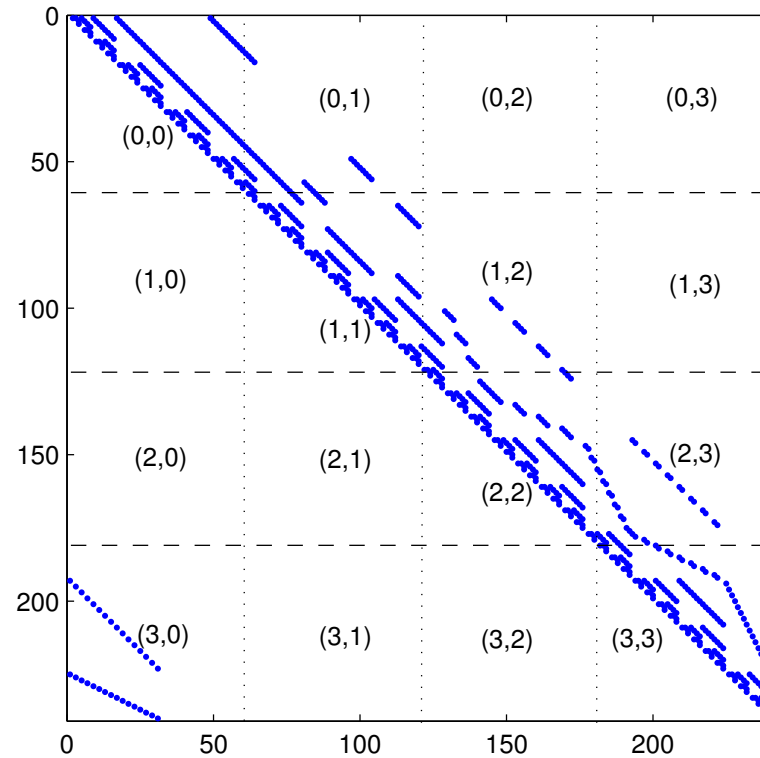
# Sparsity Pattern (Kanban and FMS)

---



# Sparsity Pattern (Polling)

---





# The Computations

---

- A revisit: MVP operation is central to
  - Krylov subspace and basic iterative methods for steady state solution
  - uniformisation method for transient solution

- An in-core MVP-based Jacobi iteration

$$r \leftarrow -Q^T \pi$$

for  $i = 0$  to  $n - 1$

$$\tilde{\pi}_i \leftarrow \pi_i + r_i / q_{ii}$$

$$r \leftarrow -Q^T \tilde{\pi}$$

Test for convergence

$$\pi \leftarrow \tilde{\pi}$$

Repeat iteration if required

# The Computations...

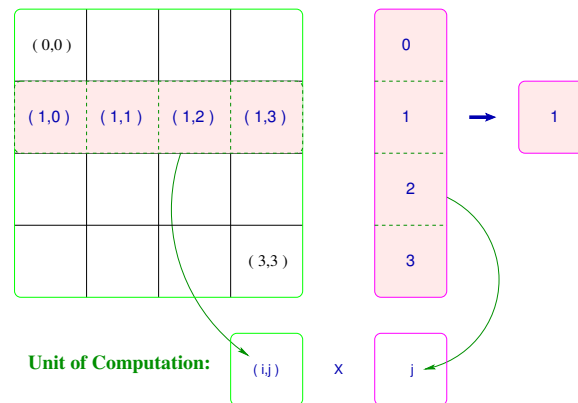
---

- Storage requirement
  - Matrix storage in some sparse format
  - Two iteration vectors:  $2 \times 8n$  Bytes
- Gauss-Seidel
  - Typically, converges faster than Jacobi
  - Requires one iteration vector
  - Standard MVP not possible!
  - Partitioning the matrix and vector into blocks can help

# An MVP-based Computation

---

- Partitioning
  - $A$ :  $B \times B$  square blocks (not necessary)
  - $x$ :  $B$  blocks with  $n/B$  entries each
- A unit of Computation: A sub-MVP
  - matrix block  $\times$  vector block ( $A_{ij} \times X_j$ )



# A (Serial) Block Jacobi Algorithm

---

`ser_block_Jac(  $\check{A}$ ,  $d$ ,  $b$ ,  $x$ ,  $P$ ,  $n[\ ]$ ,  $\varepsilon$  ) {`

1. `var  $\tilde{x}$ ,  $Y$ ,  $k \leftarrow 0$ ,  $\text{error} \leftarrow 1.0$ ,  $i$ ,  $j$ ,  $p$`
2. `while(  $\text{error} > \varepsilon$ )`
3.      `$k \leftarrow k + 1$`
4.     `for(  $0 \leq i < P$ )`
5.          `$Y \leftarrow B_i$`
6.         `for(  $0 \leq j < P$ )`
7.              `$Y \leftarrow Y - \check{A}_{ij} X_j^{(k-1)}$`
8.         `for(  $0 \leq p < n[i]$ )`
9.              `$X_i^{(k)}[p] \leftarrow D_i[p]^{-1} Y[p]$`
10.     `compute  $\text{error}$`
11.      `$X_i^{(k-1)} \leftarrow X_i^{(k)}$                      }`

# An Out-of-Core Algorithm

---

Integer constant:  $B$  (*number of blocks*)

Semaphores:  $S_1, S_2$ : occupied

Shared variables:  $R_0, R_1$  (*To read matrix  $A$  blocks into RAM*)

## Disk-IO process

1. Local variables:  $i, j, t = 0$
2. **while** not converged
3.     **for**  $i \leftarrow 0$  to  $B - 1$
4.         **for**  $j \leftarrow 0$  to  $B - 1$
5.             **if** not an *empty* block
6.                 **disk\_read** ( $A_{ij}, R_t$ )
7.             **Signal**( $S_1$ )
8.             **Wait**( $S_2$ )
9.              $t = (t + 1) \bmod 2$

## Compute process

- Local variables:  $i, j, t = 0$
- while** not converged
- for**  $i \leftarrow 0$  to  $B - 1$
- for**  $j \leftarrow 0$  to  $B - 1$
- Wait**( $S_1$ )
- Signal**( $S_2$ )
- if**  $j \neq B - 1$
- if** not an *empty* block
- sub-MVP**( $A_{ij}X_j, R_t$ )
- else**
- Update  $X_i$
- check for convergence
- $t = (t + 1) \bmod 2$

# Out-of-Core on a Single Machine

---

Model	$k$	States ( $n$ )	$a/n$	Times		Iterations
				Per iteration (seconds)	Total (hr:min:sec)	
FMS	11	54,682,992	9.5	51.6	23:16:39	1624
	12	111,414,940	9.7	170	84:54:20	1798
	13	216 427 680	9.9	327	179:34:39	1977
	14	403,259,040	10.03	1984	–	> 50
	15	724,284,864	10.18	6312	–	> 50
Kanban System	7	41,644,800	10.8	18.9	4:12:38	802
	8	133,865,325	11.3	139	38:34:21	999
	9	384,392,800	11.6	407	136:54:37	1211
	10	1,005,927,208	11.97	1424	566:49:52	1433
Polling System	22	138,412,032	11.8	143	1:28:11	37
	23	289,406,976	12.3	264	2:47:12	38
	24	603,979,776	12.8	460	4:51:20	38
	25	1,258,291,200	13.3	1226	13.16:54	39

# A Parallel Jacobi Algorithm for Process $p$

```
par_block_Jac(  $\check{A}_p, D_p, B_p, X_p, T, N_p, \varepsilon$  ) {  
  1. var  $\tilde{X}_p, Z, k \leftarrow 0, \text{error} \leftarrow 1.0, q, h, i$   
  2. while(  $\text{error} > \varepsilon$  )  
  3.    $k \leftarrow k + 1; h \leftarrow 0$   
  4.   for(  $0 \leq q < T; q \neq p$  )  
  5.     if(  $\check{A}_{pq} \neq \mathbf{0}$  )  
  6.       send(  $\text{request}_{X_q}, q$  );  $h \leftarrow h + 1$   
  7.    $Z \leftarrow B_p - \check{A}_{pp} X_p^{(k-1)}$   
  8.   while(  $h > 0$  )  
  9.     if( probe(  $\text{message}$  ) )  
 10.      if(  $\text{message} = \text{request}_{X_p}$  )  
 11.        send(  $X_p, q$  )  
 12.      else  
 13.        receive(  $X_q, q$  );  $h \leftarrow h - 1$   
 14.         $Z \leftarrow Z - \check{A}_{pq} X_q^{(k-1)}$   
 15.      serve(  $X_p, \text{request}_{X_p}$  )  
 16.      for(  $0 \leq i < N_p$  )  
 17.         $X_p^{(k)}[i] \leftarrow D_p[i]^{-1} Z[i]$   
 18.      compute error collectively }  
}
```

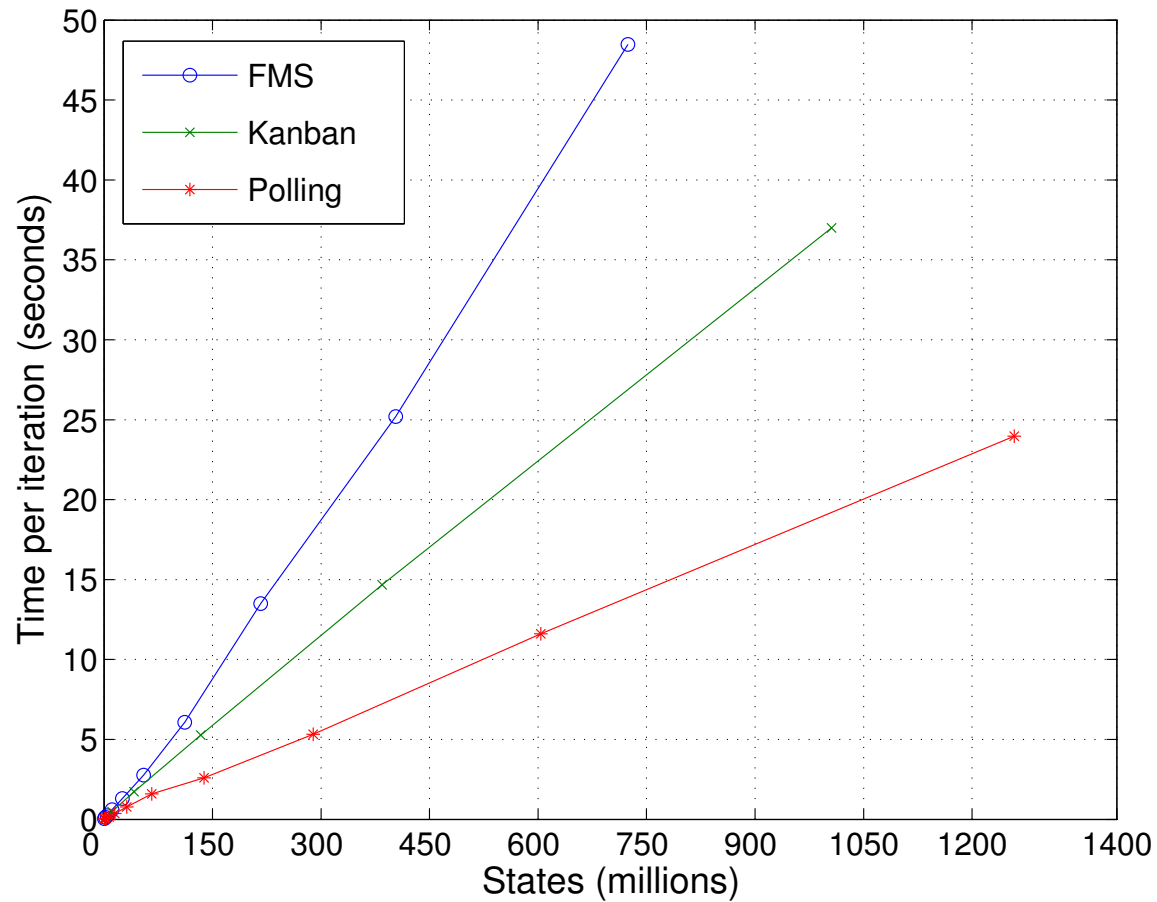
# Parallel Execution Results on 24 Nodes

$k$	States ( $n$ )	MB/Node	Time		Total Iterations
			Iteration (seconds)	Total (hr:min:sec)	
<b>FMS Model</b>					
12	111,414,940	170	6.07	3:40:57	2184
13	216 427 680	306	13.50	8:55:17	2379
14	403,259,040	538	25.20	18:02:45	2578
15	724,284,864	1137	48.47	37:26:35	2781
<b>Kanban System</b>					
7	41,644,800	53	1.73	33:07	1148
8	133,865,325	266	5.27	2:02:06	1430
9	384,392,800	564	14.67	7:03:29	1732
10	1,005,927,208	1067	37.00	21:04:10	2050
<b>Polling System</b>					
22	138,412,032	328	2.60	44:31	1027
23	289,406,976	667	5.33	1:36:02	1081
24	603,979,776	811	11.60	3:39:38	1136
25	1,258,291,200	1196	23.97	7:54:25	1190



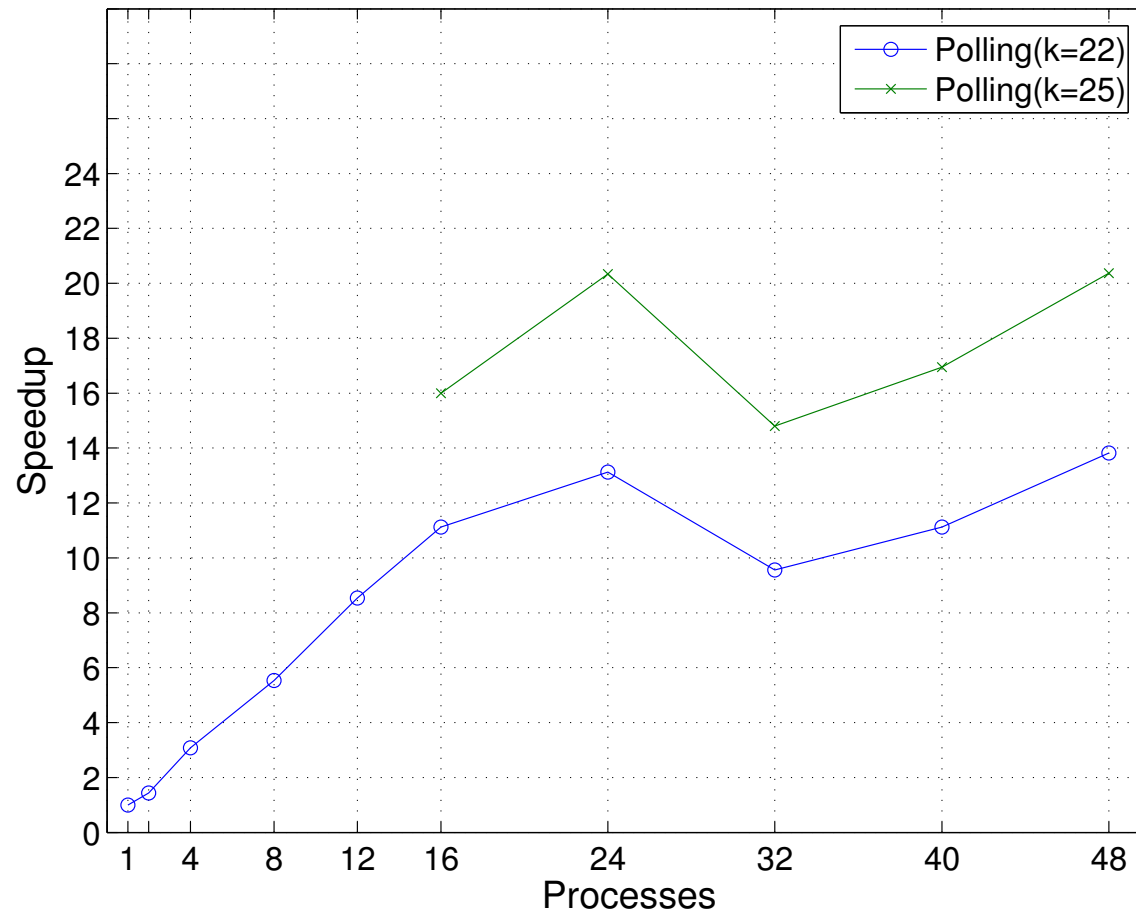
# Comparison of Solution Times (Iteration)

---



# Speedups for the Polling System

---



# Applications I am looking at

---

- Debugging and managing distributed systems
  - event-based middleware
  - modelling program behaviour
  - integration of traditional debugging and model checking
  - scheduling and optimisation
  - systems security
    - intrusion detection through analysis
- Road traffic and networks
  - real data from M4 motorway
  - city and intercity traffic

# Applications I am looking at...

---

- Optical and storage area networks
  - WDM metro area ring networks
  - SAN extensions and mirroring strategies
- Mobile ad hoc and sensor networks
  - mobility, applications, network layer...
- Communications traffic modelling
  - multimedia and services
  - transport layer protocols

# Summary and Future Work

---

- State Space Explosion
  - discrete state models: CTMCs, DTMCs, MDPs
  - symbolic storage does help
  - out-of-core and parallel solutions
  - the largest models solved, over a billion states
- Apply the developed solution techniques to
  - more interesting real-life case studies
  - build realistic models
  - gain novel insight
- Parallel Algorithms
  - make it more adaptive
  - extend to Grids and ad hoc environments